### **CYK** parsing

LING83600: Language Technology

#### The task

*Constituency* (or *constituent*) *parsing* refers to recovering the derivational structure of a sentence, usually assuming it was generated by a context-free grammar.

Shieber (1986) shows, conclusively, that human language is not even weakly context-free; and one can parse using *mildly context-sensitive grammars* (MCSGs), such as

- lexicalized tree-adjoining grammars (LTAGs) or
- combinatory categorial grammar (CCGs).

CFGs can be parsed in  $O(n^3)$  or better; MCSGs can be parsed in  $O(n^6)$ .

#### Ambiguity

Non-trivial CFGs tend to have substantial ambiguities, so we use:

- *Probabilistic* productions, making the grammar a PCFG and parsing a special case of the Viterbi algorithm
- Learned (e.g., Klein & Manning 2003) or induced (e.g., Petrov et al. 2006) enrichments to the productions to encode argument structure of verbs, etc.
- Optionally, discriminative reranking (e.g., Collins 1999, Bikel 2004)

None of these has a major impact on the actual parsing algorithm, however.

#### Structural ambiguity: prepositional phrase attachment

Pope Francis **on** Saturday appointed a victim **of** sexual abuse and a senior cardinal known **for** his zero-tolerance approach **to** a new group charged **with** advising the Catholic Church **on** how to respond **to** the problem **of** sexual abuse **of** children. (Wall St. Journal, 2014-03-22)

- The prepositional phrase on Saturday is construed as a modifier of Pope Francis rather than of appointed.
- The phrase to a new group charged with advising the Catholic Church on how to respond to the problem of sexual abuse of children is construed as a modifier of zero tolerance approach rather than of appointed.

#### CFGs and compilers

CFG parsing has been studied *extensively* for compiler design long before it was used to parse natural language.

Computer languages are designed to be unambiguous or even deterministic, and compiler designers can take advantages of more-efficient algorithms for these special cases (e.g., the Earley algorithm, the LL and LALR parsers, etc.).

Probably the best known parsing algorithm for natural language is the <u>simultaneously-invented</u> Cocke-Younger-Kasami (CYK or CKY) algorithm, a simple bottom-up parsing algorithm.



Aho & Ullman 1977





#### Aho, Lam, Sethi & Ullman 2006

Aho, Sethi & Ullman 1986

#### Example grammar

- S NP VP  $\rightarrow$ VP VP PP  $\rightarrow$ VP V NP  $\rightarrow$ VP V  $\rightarrow$ ΡP  $\rightarrow$  P NP NP  $\rightarrow$  D N NP  $\rightarrow$  NP PP NΡ N  $\rightarrow$ Pagliaccio NΡ  $\rightarrow$ V  $\rightarrow$ eats spaghetti | fork Ν  $\rightarrow$ Ρ with  $\rightarrow$ 
  - $D \rightarrow a$

- # Permits PPs to adjoin to VP.
- # Permits transitive VPs.
- # Permits intransitive VPs.
- # Permits article-noun NPs.
- # Permits PPs to adjoin to NP.
- # Permits bare NPs.

#### Eliminate unary productions

Before:

 $\begin{array}{ccc} VP & \rightarrow & V \\ NP & \rightarrow & N \end{array}$ 

After:

 $VP \rightarrow eats$ NP  $\rightarrow$  spaghetti, fork

#### Grammar representation structure

The preterminal rules are simple mappings, and only used once.

```
preterminals: Dict[str, List[str]]
```

The non-terminal rules map from pairs of strings (the children/daughters/productions) to non-terminals.

```
nonterminals: Dict[Tuple[str, str], List[str]]
```

Looking up productions to find candidate non-terminals is known as grammar intersection. Intersection is performed  $O(n^3)$  times, so it has to be constant-time.

## Pagliaccio eats spaghetti with a fork

(Do you see the ambiguity?)



#### The ambiguity

1. Pagliaccio  $[_{VP} [_{VP} eats spaghetti]$  [with a fork] ]

(The fork is an instrument for eating.)

2. Pagliaccio eats [<sub>NP</sub> [<sub>NP</sub> spaghetti] [with a fork] ]

(The fork is a tasty topping for the spaghetti.)

#### The CYK chart

Given a sentence of length *n*, construct an *n* x *n* table, empty above the diagonal, in which each cell is a container of non-terminal labels.

For efficient memory use you **might** want to construct a Chart class that has a *n* x *n* / 2 + 1 array and overloads \_\_getitem\_\_ and \_\_setitem\_\_ so you can pretend you really do have a *n* x *n* table without wasting the extra space.

# (If you find an error in the demo, let me know.)

#### Initial chart (plus an extra row at the bottom, for display)



#### Insert preterminals

NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork









































		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

(This is the NP-adjunction analysis.)

		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork









		_			
	VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

(This is the NP-adjunction analysis.)

		_			
	VP, <b>VP</b>				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

(This is the VP-adjunction analysis.)

		_			
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

S, S		_			
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

S, S					
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

S, S					
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

S, S					
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

S, S					
	VP, VP				
		NP			
S			PP		
S	VP			NP	
NP	VP	NP	Р	D	NP
Pagliaccio	eats	spaghetti	with	а	fork

#### Recognition vs. parsing

When there an S in the top-left cell of the chart, the sentence is *recognized* by the grammar; i.e., there is at least one complete parse of the sentence.

To actually *parse* (and to extract the parses), we have to keep backtraces similar to Viterbi decoding.

#### Generalizing to PCFGs

Whereas above each cell was a container of non-terminals, in PCFG parsing each cell is a container of triples of;

- 1. non-terminal
- 2. log-probability
- 3. backtraces

Much like with the Viterbi algorithm, we obtain the **highest-probability parse** by tracing back from the **highest-probability S-triple in the top-left cell**.

#### Data

Unlike dependency parsing (the subject of next week's lecture), constituency parsing treebanks are available for maybe a dozen languages, and many are proprietary. Two widely studied ones are:

- English: The Wall St. Journal portion of the Penn Treebank, 3rd edition (Marcus et al. 1999; first release was 1993, I think)
- German: The Tiger2 corpus (Brants et al. 2002)

For German, which is "morphologically rich and less-configurational", it is generally agreed that parsing is poor without the assistance of a high-quality morphological analyzer (e.g., Dubey 2005, Fraser et al. 2013, Dehouck and Denis 2018).

#### Evaluation

Parsing performance is usually evaluated using the *F*-score of the constituents:

- Precision: the number of the constituents in the candidate/hypothesis parse that are present in the gold/reference parse
- Recall: the number of constituents in the gold/reference parse that are present in the candidate/hypothesis parse
- F-score: the harmonic mean of precision and recall

PARSEVAL (Black et al. 1991), a widely used evaluation script, also includes a measure that counts *crossing brackets*.

#### Software

- Charniak (2000) is the best-known of the PCFG-plus-discriminative-reranking parsers and is still used for feature extraction (etc.)
- The BUBS "grammar-agnostic" parser (Bodenstab et al. 2011, Dunlop et al. 2011) is a good tradeoff between simplicity and generality, speed and accuracy; unfortunately it's in Java

#### This week

Please read:

- Eisenstein ch. 10, and
- (optionally) Jurafsky & Martin ch. 12-14.

#### Project ideas

- Implement either of the "Berkeley" parsers (Klein & Manning or Petrov et al.) and evaluate using the standard split of the Wall St. Journal portion of PTB-3.
- Perform a hybrid (automatic-manual) error analysis of prepositional phrase attachment ambiguity.
- Study the effects of morphological features in constituency parsing along the lines of Fraser et al. 2013, etc.